Symbolic Execution of Security Protocol Impl.: Handling Cryptographic Primitives

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WiFi®

4-way handshake









Handling Crypto



4-way handshake



return

```
int num = len/data[2]
```

Symbolic Execution

```
data[0] != 1
void recv(data, len) {
  if (data[0] != 1)
    return
  if (data[1] != len)
    return
  int num = len/data[2]
```

data[0] == 1
<pre>void recv(data, len) {</pre>
if (data[0] != 1)
return
<pre>if (data[1] != len)</pre>
return
<pre>int num = len/data[2]</pre>
• • •
}





Symbolic Execution

```
data[0] == 1 &&
     data[1] == len
void recv(data, len) {
  if (data[0] != 1)
    return
  if (data[1] != len)
    return
                           Can data[2] equal zero
  int num = len/data[2]
                            under the current PC?
```

Symbolic Execution

```
data[0] == 1 &&
     data[1] == len
void recv(data, len) {
  if (data[0] != 1)
                              Yes! Bug detected!
    return
  if (data[1] != len)
    return
                           Can data[2] equal zero
  int num = len/data[2]
                           under the current PC?
```

Implementations

> Works on LLVM bytecode

We build upon KLEE

- Actively maintained

Practical limitations:

- $|paths| = 2^{|if-statements|}$
- Infinite-length paths
- > SMT query complexity









4-way handshake



Motivating Example
 Mark data as symbolic
void recv(data, len) {

- plain = decrypt(data, len)
 - if (plain == NULL) return

if (plain[0] == COMMAND) process_command(plain) else

• • •

Motivating Example Mark data as symbolic void recv(data, len) { Summarize crypto algo. plain = decrypt(data, len) ← (time consuming) if (plain == NULL) return Analyze crypto algo. if (plain[0] == COMMAND) + (time consuming) process command(plain) else Won't reach this code!

Efficiently handling decryption?

Decrypted output

fresh symbolic variable

Example Mark data as symbolic void recv(data, len) { plain = decrypt(data, len) { create fresh if (plain == NULL) return symbolic variable symbolic variabl

if (plain[0] == COMMAND) process_command(plain) Normal analysis else

... → Can now analyze code that parses decrypted data

Other Applications

Handling hash functions

- > Output = fresh symbolic variable
- > Also works for HMACs (Message Authentication Codes)



Tracking use of crypto primitives?

- > Recording relationship between input & output
- > Treating fresh variable as information flow taint

Detecting Crypto Misuse



Timing side-channels

- > \forall (*paths*): all bytes of MAC in path constraint?
- > If not: comparison exits on first difference



Decryption oracles

- > Behavior depends on unauth. decrypted data
- > Decrypt data is in path constraint, but not in MAC







Handling Crypto



4-way handshake





Used to connect to any protected Wi-Fi network



Mutual authentication



Negotiates fresh PTK: pairwise transient key

4-way handshake (simplified)

 \leftarrow - - - - - optional 802.1x authentication - - - - - \rightarrow



















Symbolic execution of







Intel's iwd deamon

wpa_supplicant

kernel driver

How to get these working under KLEE?





Avoid running full program under KLEE

> Would need to model Wi-Fi stack symbolically

Our approach

- > iwd contains unit test for the 4-way handshake
- > Reuse initialization code of unit test!
- > Symbolically execute only receive function

wpa_supplicant



Unit test uses virtual hardware and runs full AP

> Still need to simulate Wi-Fi stack...

Alternative approach:

- > Write unit test that isolates 4-way handshake like iwd
- > Then symbolically execute receive function!
- > Need to modify code of wpa_supplicant (non-trivial)

MediaTek's Driver



- No unit tests & it's a Linux driver
- > Symbolically executing the Linux kernel?!

Inspired by previous cases

- > Write unit test & simulate used kernel functions in userspace
- > Verify extracted code is correctly simulated in userspace!

Not all our unit tests are created equally



https://github.com/vanhoefm/woot2018







Handling Crypto



4-way handshake



Discovered Bugs I



Timing side-channels

- > Authentication tag not checked in constant time
- > MediaTek and iwd are vulnerable



Denial-of-service in iwd

- > Caused by integer underflow
- > Leads to huge malloc that fails

Discovered Bugs II



Buffer overflow in MediaTek kernel driver
Occurs when copying the group key
May lead to remote code execution



Flawed AES unwrap crypto primitive
Also in MediaTek's kernel driver
Manually discovered

Decryption oracle in wpa_supplicant



Decryption oracle:

- Doesn't check authenticity of malformed handshake message
- > But does decrypt and process data

→ Decrypt group key (GTK) in Message 3 (Msg3)

Decryption oracle in wpa_supplicant II

Msg3': decrypted using RC4, but not authenticated





 \rightarrow Parsing only succeeds if x_{37} is zero

Future work

Short-term

- > Efficiently simulate reception of multiple packets
- > If 1st packet doesn't affect state, stop exploring this path

Long-term

- > Extract packet formats and state machine
- > Verify basic properties of protocol

Conclusion



- > Symbolic execution of protocols
- Simple simulation of crypto
- > Interesting future work

As a final note...

I wrote a vulnerability scanner that abstracts all the predicates in a binary, traverses the callgraph and generates phormulaes to run then with a SMT solver. I found 1 vuln in 3 days with this tool.

He wrote a dumb ass fuzzer and found 5 vulns in 1 day.

Good thing I'm not a n00b like that guy.



Thank you!

Questions?