

# Operating Channel Validation

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# Contributions



} Paper: attacks & high-level defense

} Specification: text for inclusion in 802.11

} Implementation: modified hostap

# Old attacks don't need Man-in-the-Middle (MitM)

```
2) 6E(38400) 81(37376) 79(36864)
0) 15(38656) 7B(38400) BB(37888)
8) 23(38144) 97(37120) 59(36608)

KEY FOUND! [ 1F:1F:1F:1F:1F ]
tly: 100%
```

Breaking WEP



Dictionary attacks



Breaking WPS



Rogue APs

# New attacks do require MitM



## Traffic Analysis

- › **Capture all** encrypted frames
- › **Block** certain encrypted frames

## Attacking broadcast TKIP

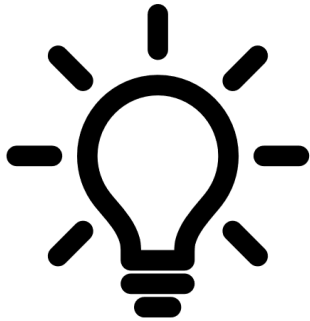
- › **Block** MIC failures
- › **Modify** encrypted frames

Chop/Chop

# New attacks do require MitM

Exploit implementation bugs

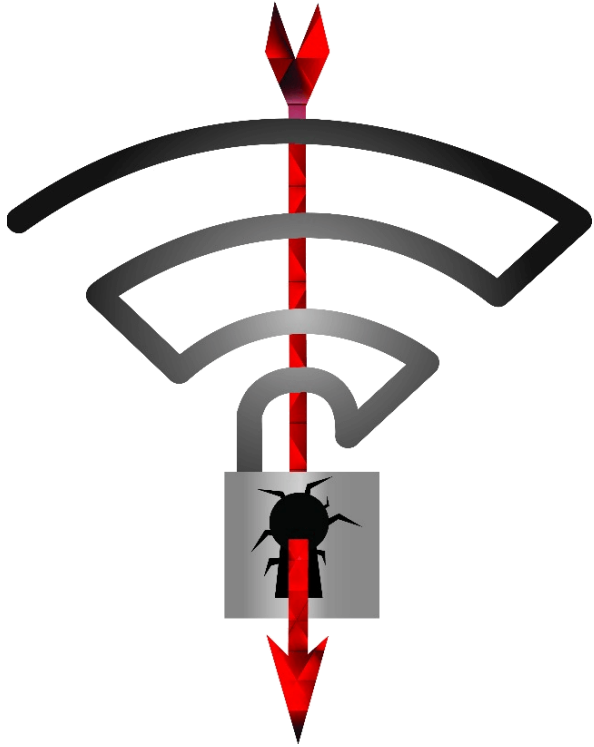
- › **Block** certain handshake messages
- › E.g. bugs in 4-way handshake



New attack scenarios

- › See paper for details
- › E.g. **modify** advertised capabilities

# The elephant in the room



## Key Reinstallation Attacks (KRACKs)

- › **Block & delay** handshake frames
- › E.g. 4-way & group handshake

## Not all KRACKs require MitM

- › E.g. FT handshake (802.11r)

# Obtaining multi-channel MitM

Clone AP on **different channel!**



# Force client on rogue channel?



Jam channel of real AP

- › Victim will connect on rogue AP
- › Stop jamming when client connects

We found an easier way while making the defense!

- › **Abuse channel switch announcements**



# Channel Switch Announcements (CSAs)

## Background:

- › AP may dynamically switch channels
- › E.g. when radar pulses are detected
- › Sends CSAs to connected clients
- › **Clients switch to new channel in CSA**

## Adversary can forge CSAs

- › **Abuse to switch victim to rogue channel!**



# Can we prevent MitMs?

## Threat model

- › Focus on verifying channel and bandwidth
- › We exclude low-layer attacks such as beamforming

**Goal is to make attacks harder, not impossible!**

Similar to the idea of stack canaries.

# Proposed Defense

Verify operating channel when connecting to a network

- › E.g. in the 4-way and FT handshake

Also verify channel in

- › WNM-Sleep exit frames: avoid tricky edge cases
- › Group key handshake: defense in depth

# Encoding the current channel

Operating Channel Information **(OCI) element:**

Operating class	Channel number	Segment index 1
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1. Operating class: defines the bandwidth
2. Channel number: defines primary channel
  - › Together this also defines the central frequency
3. Seg idx 1: for 80+80 MHz channels

# Problem: Channel Switch Announcements (CSAs)

## Unauthenticated CSAs

- › Need to verify securely

## Authenticated CSAs

- › May not arrive → need to verify reception!

**Solution: authenticate CSA using SA query**

# Limitations

Other (partial) MitM attacks still possible:

- › Partial MitM when client didn't receive CSA
- › Adversary can act as repeater
- › Other physical-layer tricks

So why use this defense?

- › **Remaining attacks are harder & not always possible**
- › Straightforward to implement

# Standardization efforts

March 2018

doc.: IEEE 802.11-17/1807r10

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## IEEE P802.11 Wireless LANs

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### Defense against multi-channel MITM attacks via Operating Channel Validation

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- › Detailed technical specification
- › **Has extra discussions not present in paper!**
- › Hopefully ratified soon 😊

# Proof-of-concept

[github.com/vanhoefm/hostap-channel-validation](https://github.com/vanhoefm/hostap-channel-validation)



- › Code for 4-way handshake
- › Other handshakes in progress

Some remarks:

- › Has many automated tests!
- › Kernel may change bandwidth



# Conclusion



- › Easy MitM with channel switches
- › We prevent multi-channel MitM
- › Other MitM still possible
- › Being standardized!

Thank you!

Questions?