Fragment and Forge: Breaking Wi-Fi Through Frame Aggregation and Fragmentation



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Advancements in Wi-Fi security

- > WPA3 is continously being updated
 - » Preventing recent Dragonblood [VR20] attack
 - » Securing hotspots using asymmetric crypto
- > Operating channel validation [VBDOP18]
- > Beacon protection [VAP20]
- > KRACK patches proven secure [CKM20]

Despite these major advacements, found **flaws in all networks** (incl. WPA2/3)

Design flaws

Design flaws

Aggregation

Mixed key

Fragment cache



Sending small frames causes high overhead:



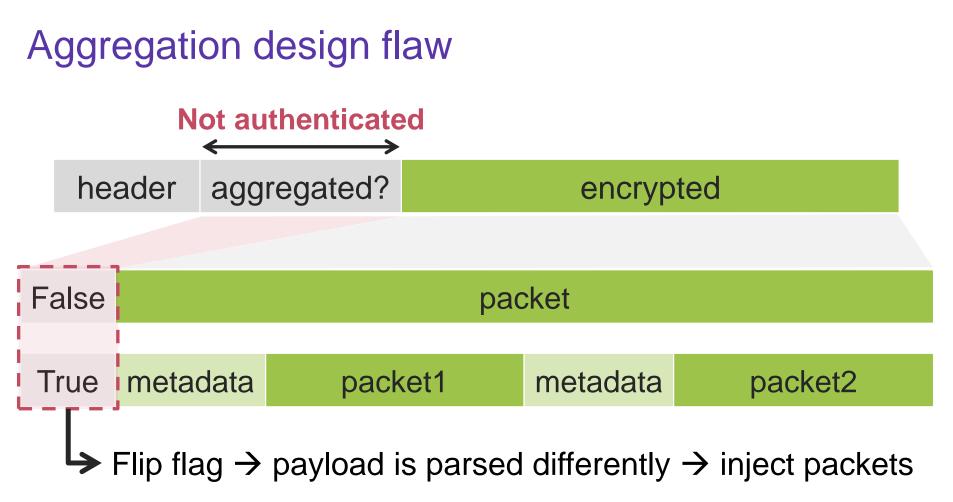
This can be avoided by **aggregating frames**:

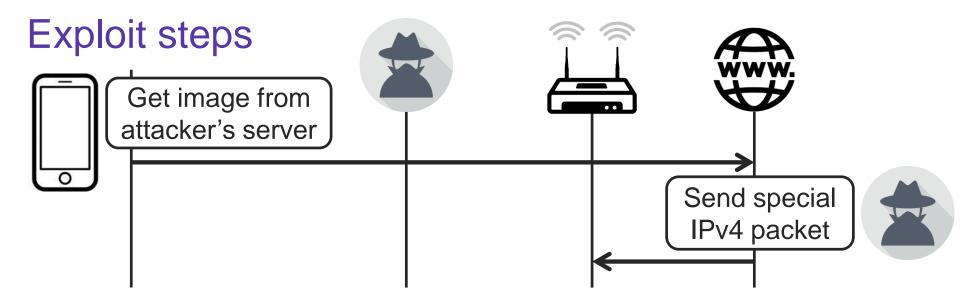
header' packet1 packet2 ... ACK

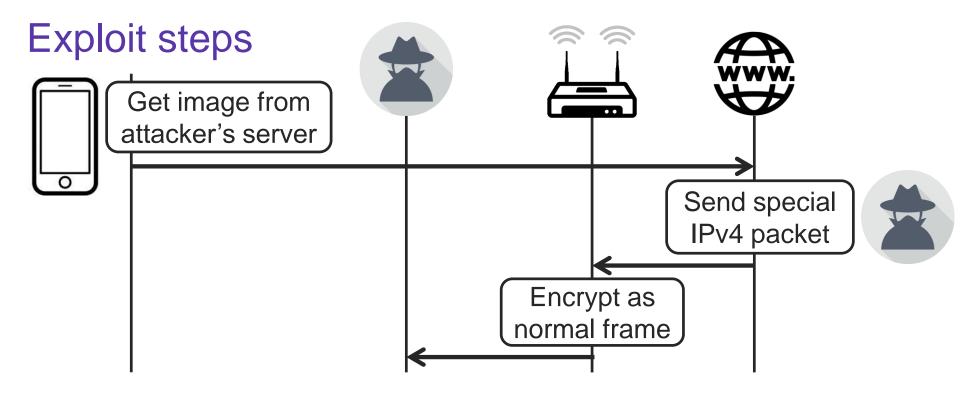
Problem: how to recognize aggregated frames?

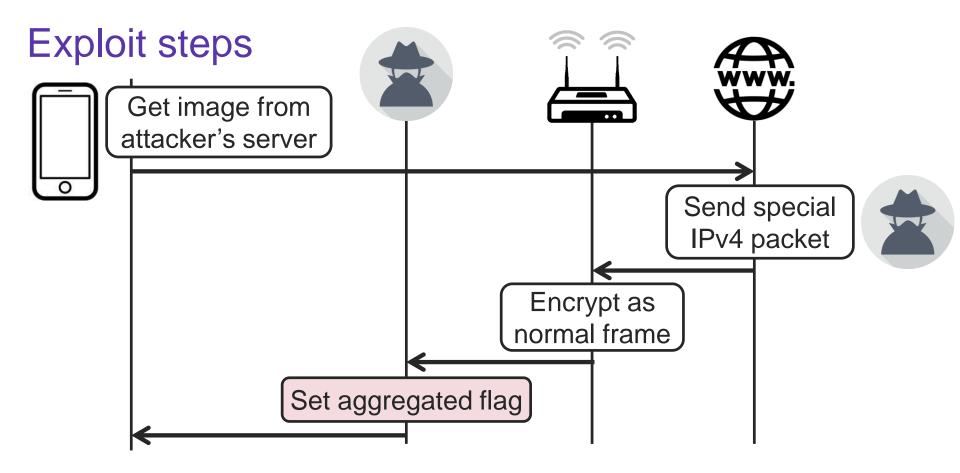
Aggregation design flaw

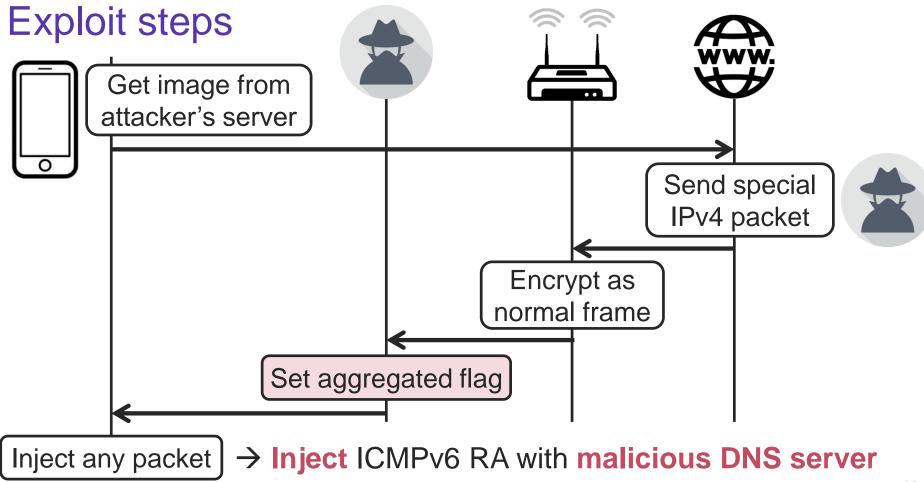
header a		aggregated? encrypted			ted	
False			packet			
True	metada	ita pack	et1	metadata	packet2	

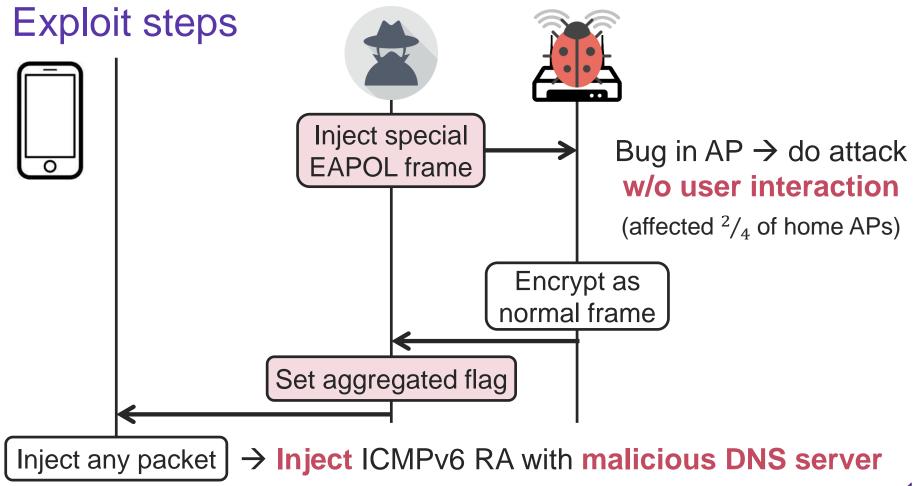












Aggregation

Implementation Flaws

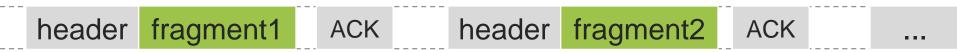
edFragmentycache



Large frames have a high chance of being corrupted:



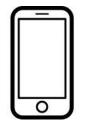
Avoid by **fragmenting** & only retransmitting lost fragments:

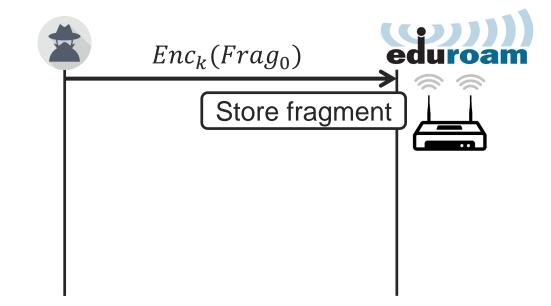


 \rightarrow Protected header info defines place in original frame

Fragment cache design flaw

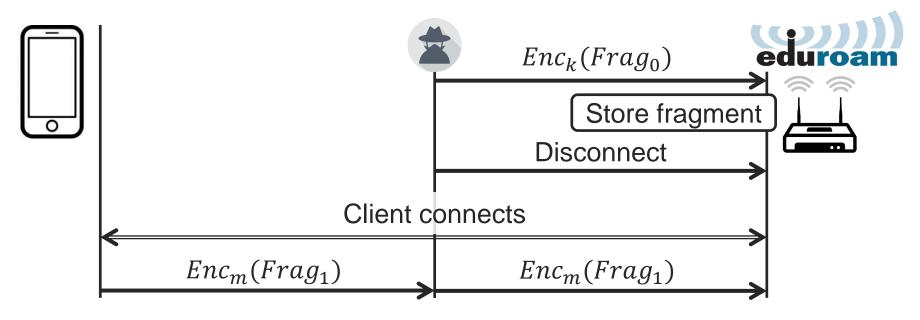
Fragments aren't removed after disconnecting:





Fragment cache design flaw

Fragments aren't removed after disconnecting:



> Attacker's $Frag_0$ and client's $Frag_1$ is reassembled

Summary of impact

Abuse to **exfiltrate or inject packets** assuming:

- 1. Hotspot-like network where users distrust each other
- 2. Client sends fragmented frames (rare unless Wi-Fi 6)
- Even the ancient **WEP protocol is affected**!
- > WEP is also affected by the mixed key design flaw

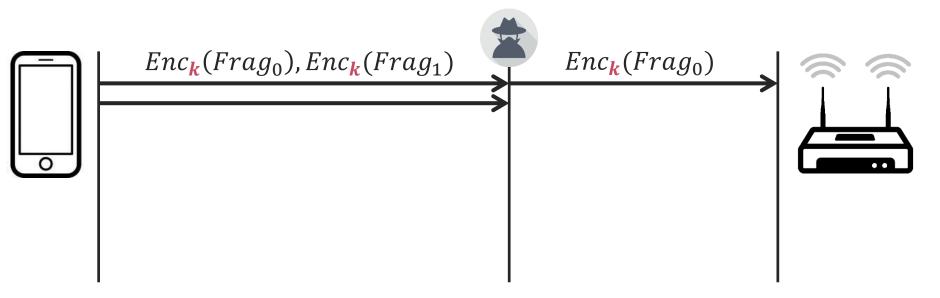
→ Design flaws have been part of Wi-Fi since 1997

Aggregation

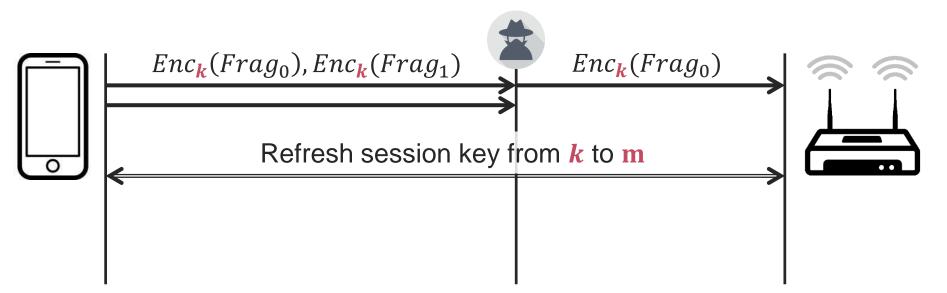
Mixed key

Fragment cache

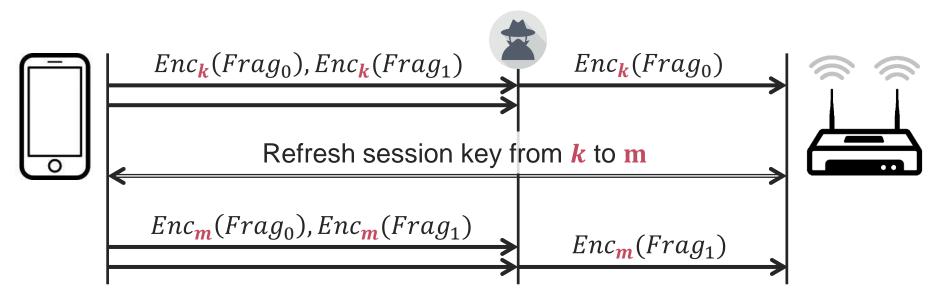
Fragments decrypted with different keys are reassembled:



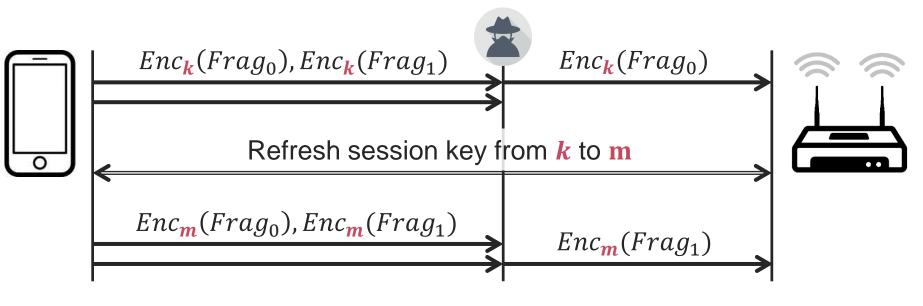
Fragments decrypted with different keys are reassembled:



Fragments decrypted with different keys are reassembled:



Fragments decrypted with different keys are reassembled:



→ Can mix fragments of different frames

Summary of impact

Abuse to **exfiltrate data** assuming:

- 1. Someone sends fragmented frames (rare unless Wi-Fi 6)
- 2. Victim will connect to server of attacker
- 3. Network periodically refreshes the session key

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Abuse to **exfiltrate data** assuming:

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 - » Combine with implementation flaw to avoid this condition

Design flaws

	Plaintext frames		
Design flaws	Broadcast fragments		
	Cloacked A-MSDUs		

Trivial frame injection

Plaintext frames wrongly accepted:

- > Depending if fragmented, broadcasted, or while connecting
- > Sometimes frames that **resemble a handshake** message
- > Examples: Apple and some Android devices, some Windows dongles, home and professional APs, and many others!

→ Can trivially inject frames

No fragmentation support

No fragmentation support

Some devices don't support fragmentation

- > But they treat fragmented frames as full frames
- > Examples: OpenBSD and Espressif chips

→ Abuse to **inject frames** under right conditions

→ All devices are vulnerable to one or more flaws

Created tool to test devices

Has 45+ test cases for both clients and APs:

Command	Short d	Non-consecutive PNs attack (§6.2)		Command	Short de	Mixed plain/encrypt attack (§6.3)	
Sanity checks		ping I,E,Einc-pn 2	Send a fragmented ping with non-	A-MSDU attacks (§3)		ping I,E,Eamsdu	Send a normal ping as a fragmente
ping	Send a normal ping.	Mixed plain/encrypt attack (§6.3)		ping I,Eamsdu-fake	If this test succeeds, the A-MSDU fl	ping I,E,P,E	Ping with first frag. encrypted, seco
ping I,E,E	Send a normal fragmented ping.	ping I,E,P	Send a fragmented ping: first fragn	ping I,Eamsdu-fakeamsdu-spp	Check if the A-MSDU flag is authen	linux-plain 3	Same as linux-plain but decoy frag
Basic device behaviour		ping i,c,r	send a fragmented ping, first fragm	ping i,camsuu-takeamsuu-spp	Check If the A-MISDU flag is authen	Broadcast checks (extensions of §6.4)	
ping I,E,Edelay 5	Send a normal fragmented ping with a	ping I,P,E	Send a fragmented ping: first fragn	Mixed key attacks (§4)			
ping-frag-sep	Send a normal fragmented ping with fra	ping I,P	Send a plaintext ping.	ping I,F,BE,E	In case the new key is installed relat	ping I,Pbcast-ra	Ping in a plaintext broadcast frame
ping-frag-seppn-per-qos	Same as above, but also works if the tar					ping BPbcast-ra [bcast-dst]	Ping in plaintext broadcast frame c
A-MSDU attacks (§3)		ping I,P,P	Send a fragmented ping: both frag	ping I,E,F,AE	Variant if no data frames are accept	ping BP [bcast-dst]	Ping in a plaintext frame during the
ping I,Eamsdu	Send a ping encapsulated in a normal (linux-plain	Mixed plaintext/encrypted fragmer	ping I,E,F,AErekey-plain	If the device performs the rekey har	eapfrag BP,BP	Experimental broadcast fragment a
amsdu-inject	Simulate attack: send A-MSDU frame w	Broadcast fragment attack (§6.4)		ping I,E,F,AErekey-plainrekey-req	Same as above, and actively reques	A-MSDU EAPOL attack (§6.5)	Experimental broadcast nagment t
amsdu-inject-bad	Same as above, but against targets that	ping I.D.Pbcast-ra	Send a unicast ping in a plaintext b	ping I,E,F,AErekey-early-install	Install the new key after sending me		
Mixed key attacks (§4)					, ,	eapol-amsdu[-bad] BPbcast-dst	Same as eapol-amsdu BP but easie
ping I,F,BE,AE	Inject two fragments encrypted under a	ping D,BPbcast-ra	Same as above, but frame is sent d	<pre>ping I,E,F,E [rekey-p1] [rekey-req]</pre>	Same as above 4 tests, but with lon	AP forwards EAPOL attack (§6.6)	
ping I,F,BE,AEpn-per-qos	Same as above, but also works if the tar	A-MSDU EAPOL attack (§6.5)		ping I,F,BE,AEfreebsd	Mixed key attack against FreeBSD c	eapol-inject 00:11:22:33:44:55	Test if AP forwards EAPOL frames k
Cache attacks (§5)		eapol-amsdu I,P	Send a plaintext A-MSDU containin	Cache attacks (§5)		eapol-inject-large 00:11:22:33:44:55	Make AP send fragmented frames
ping I,E,R,AE	Inject a fragment, try triggering a reasso	eapol-amsdu BP	Same as above, but the frame is see	<pre>ping I,E,R,AEfreebsd [full-reconnect]</pre>	Cache attack specific to FreeBSD im	No fragmentation support attack (§6.8)	
ping I,E,R,E	Same as above, but with a longer delay					1.5 11 1.7	
ping I,E,R,AEfull-reconnect	Inject a fragment, deauthenticate and re	eapol-amsdu-bad I,P	Send malformed plain. A-MSDU co	<pre>ping I,E,R,APfreebsd [full-reconnect]</pre>	Cache attack specific to FreeBSD im	ping I,D,E	Send ping inside an encrypted sec
ping I,E,R,Efull-reconnect	Same as above, but with a longer delay	eapol-amsdu-bad BP	Same as above, but the frame is see	<pre>ping I,E,R,AP [full-reconnect]</pre>	Cache attack test where 2nd fragme	ping I,E,D	Send ping inside an encrypted first

 \rightarrow Available at https://github.com/vanhoefm/fragattack

Discussion

Design flaws took two decades to discover

- > Without modified drivers some attacks will fail
- > Fragmentation & aggregation wasn't considered important

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Long-term lessons:

- > Adopt defences early even if concerns are theoretic
- > Isolate security contexts (data decrypted with different keys)
- > Keep fuzzing devices. Wi-Fi Alliance can help here!

Conclusion



> Discovered three design flaws

> Multiple implementation flaws

> Several flaws are trivial to exploit
> More info: www.fragattacks.com